## Alice and Bob in Cipherspace - Brian Hayes

**Computing arbitrary functions of encrypted data - Craig Gentry** 

### Secure Information Aggregation for Smart Grids Using Homomorphic Encryption - Fengjun Li, Bo Luo, Peng Liu

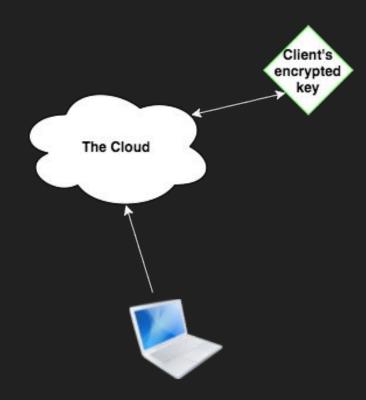
Michael Dubell & Emelie Widegren

## **Motivation**/**Problem**

## Live demo



## **Current state of Privacy**



## Can *you* trust the cloud?

#### What if the cloud changes its code?

## How does homomorphic encryption solve privacy?



• Only user can decrypt

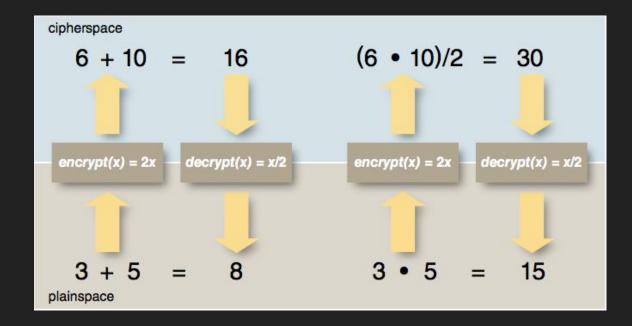


- Cloud can only do mathematical operations on the ciphertext
- Can not decrypt ciphertext



## Tell us more about this crypto magic

## **Partial homomorphic encryption**



## **El Gamal**

## Benaloh

## RSA

## Naccache-Stern

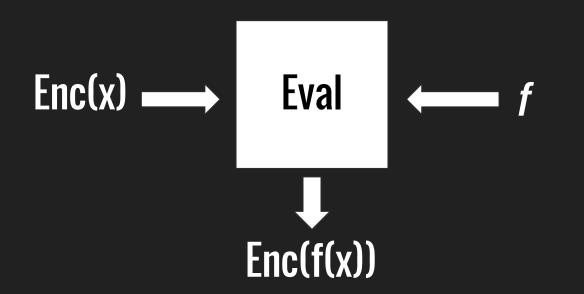
## **Boneh-Goh-Nissim**

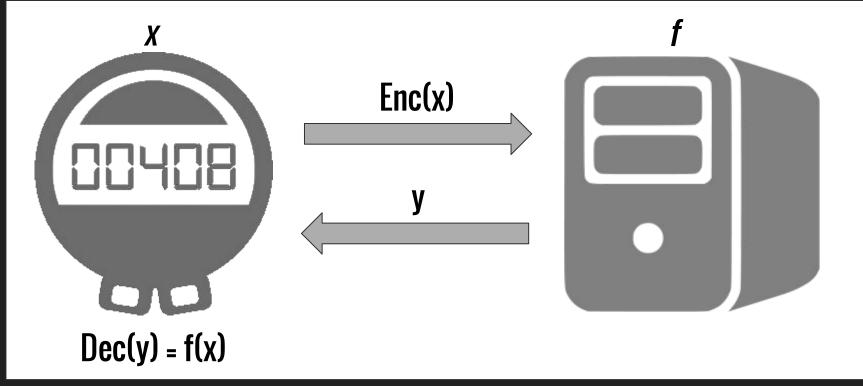
Paillier

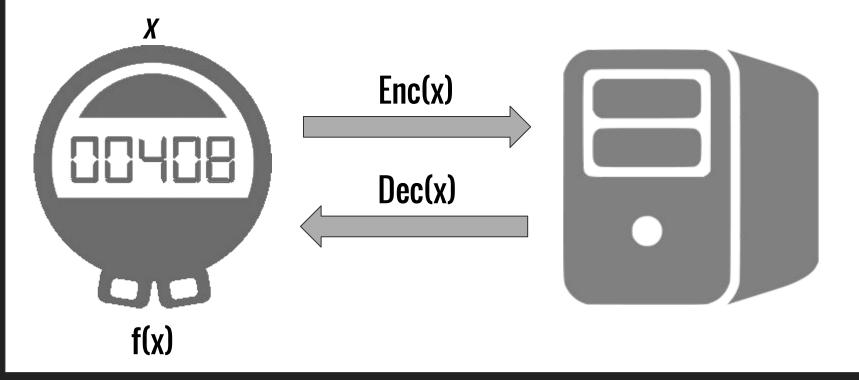
## Goldwasser-Micali

## **Fully homomorphic encryption**

- Supports arbitrary computation on ciphertexts







## **Fully homomorphic encryption**

# From a Somewhat Homomorphic Encryption scheme to a Fully Homomorphic Encryption scheme

**Craig Gentry** 

## An Encryption Scheme:

KeyGen<sub>ε</sub>(λ) -> (sk, pk) Encrypt<sub>ε</sub>(pk,m) -> c

 $Decrypt_{\epsilon}(sk,c) \rightarrow m'$ 

 $Evaluate_{\epsilon}(f,c_{1},...,c_{i}) \rightarrow c$ 

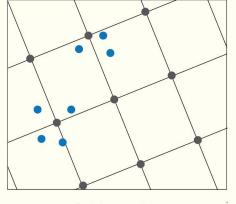
The secret key, sk, is a random P-bit odd integer p.

# A somewhat homomorphic encryption scheme

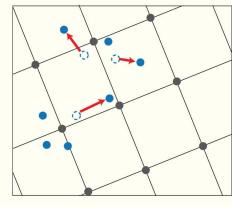
- A limited number of functions supported.

## Lattice-based cryptography

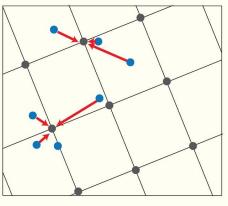
#### Lattice-based cryptography (for fully homomorphic encryption)



1. Encrypt



2. Add noise





## Bootstrapping

noise = 
$$p/2$$

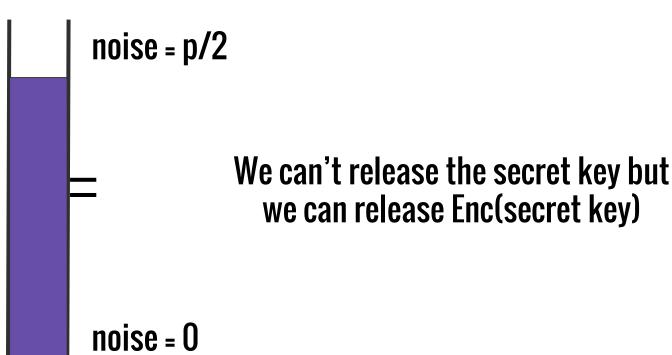
### Addition doubles noise Multiplication squares

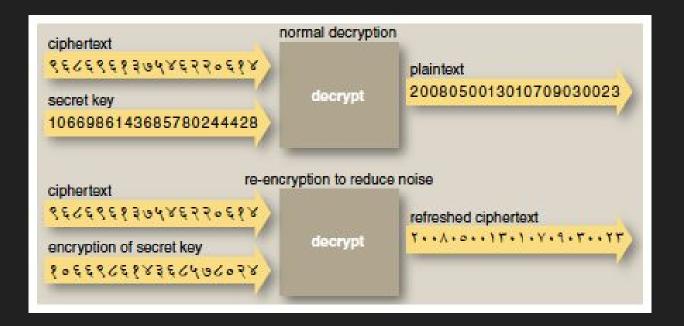
noise = O

## Bootstrapping



## Bootstrapping

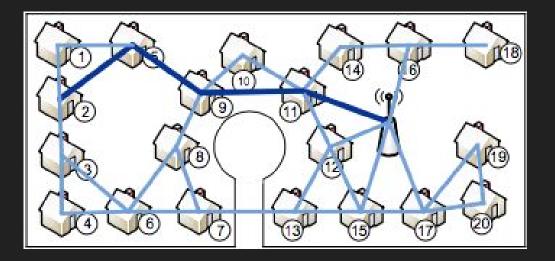




## Is Gentry's scheme practical?

# Implementing homomorphic encryption in the smart grid

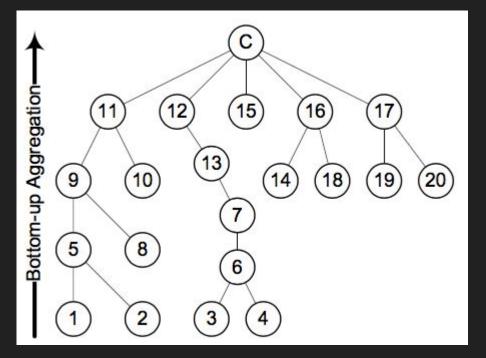
# An example of smart grid communication in a neighborhood.



[-] Excessive network traffic

[-] High overhead for collector

## Improved communication graph

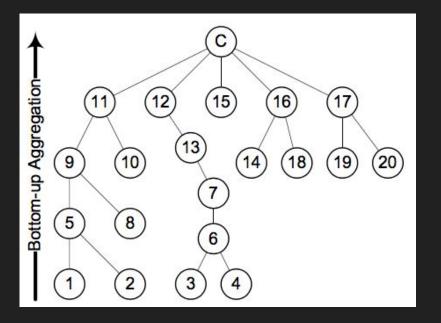


Collector calculates a BFS Tree [+] Collector offloads computation [+] Computation is done in a distributed manner

+] Tree can easily be rebalanced

# How is homomorphic encryption used in this scheme?

## System design



#### 7-tuple message

{TID, Trigger, Data, Collect, Operation, Destination, Key}

#### Child node

Generate its readings, encrypts data and sends to parent

#### **Parent node**

Collects data from all children, appends its own reading Sends data to its parent or collector

#### **Collector node**

- Collects data from children
- Performs some operation
- Decrypts data

## Paillier Cryptosystem - Key Generation

#### **Key Generation**

- 1. Pick two large prime numbers p and q;
- 2. N = p  $\cdot$  q and  $\lambda$  = lcm(p-1, q-1), where lcm represents least common multiple.
- 3. Select a random number g where g  $\in$   $Z^*_{n^2}$
- 4. Set function L(u) as: L(u) = (u 1)/N

- 5. Ensure that N divides the order of g: check if  $L(g^{\lambda} \mod N^2)$  and n are co-prime, i.e.  $gcd(L(g^{\lambda} \mod N^2), N)=1$
- 6. (N, g) is the public key pair.
- 7. (p, q) is the private key pair.

## Paillier Cryptosystem - Encryption

- 1. We want to encrypt the message:  $m \in Z_{N}^{*}$
- 2. Select a random number:  $r \in Z_{N}^{*}$
- 3. Encrypt m using:  $c = E(m) = g^{m} \cdot r^{N} \mod N^{2}$

## Paillier Cryptosystem - Decryption

1. We want to decrypt ciphertext:  $c \in Z^*_{N^2}$ 

2) Decrypt with:  $m = \mathsf{D}(c) = \left(\frac{L(c^{\lambda} \mod N^2)}{L(g^{\lambda} \mod N^2)}\right) \mod N$ 

Given 
$$c_1 = E(m_1)$$
 and  $c_2 = E(m_2)$ ,  $\forall m_1, m_2 \in Z_N$ , we have  $D(c_1 \cdot c_2 \mod N^2) = m_1 + m_2 \mod N$ 

### **Medical Applications**

### **Financial Applications**

# Advertising and Pricing

Voting

## Summary

- Properties of homomorphic encryption
- Partial homomorphism
- Fully homomorphism
- An example how it can be integrated into the smart grid