Alice and Bob in Cipherspace - Brian Hayes

Computing arbitrary functions of encrypted data - Craig Gentry

Secure Information Aggregation for Smart Grids Using Homomorphic Encryption - Fengjun Li, Bo Luo, Peng Liu

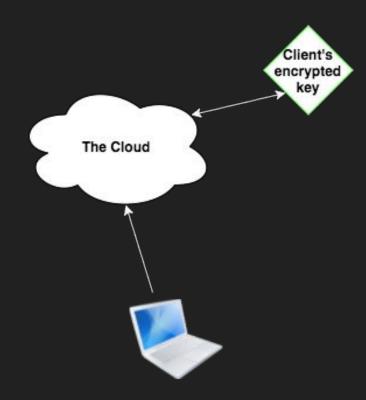
Michael Dubell & Emelie Widegren

Motivation/**Problem**

Live demo



Current state of Privacy



Can *you* trust the cloud?

What if the cloud changes its code?

How does homomorphic encryption solve privacy?



• Only user can decrypt

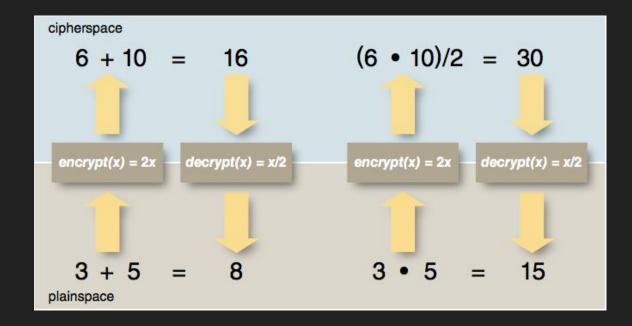


- Cloud can only do mathematical operations on the ciphertext
- Can not decrypt ciphertext



Tell us more about this crypto magic

Partial homomorphic encryption



El Gamal

Benaloh

RSA

Naccache-Stern

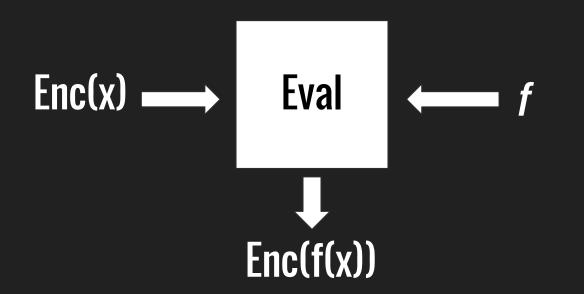
Boneh-Goh-Nissim

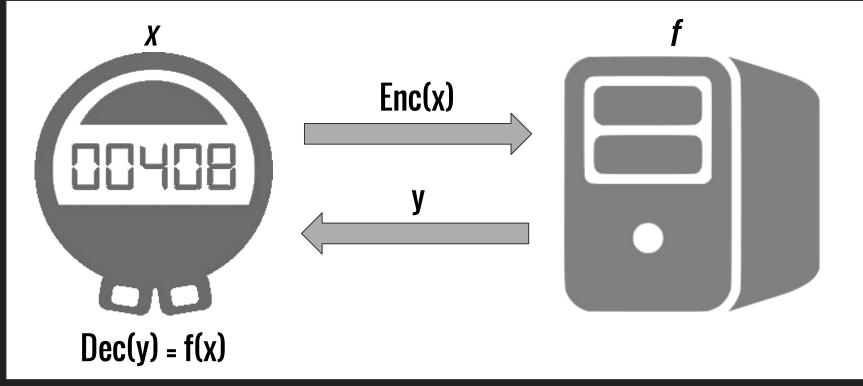
Paillier

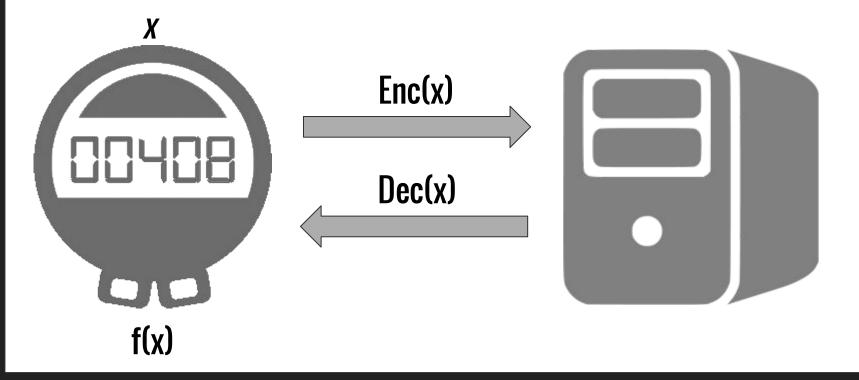
Goldwasser-Micali

Fully homomorphic encryption

- Supports arbitrary computation on ciphertexts







Fully homomorphic encryption

From a Somewhat Homomorphic Encryption scheme to a Fully Homomorphic Encryption scheme

Craig Gentry

An Encryption Scheme:

KeyGen_ε(λ) -> (sk, pk) Encrypt_ε(pk,m) -> c

 $Decrypt_{\epsilon}(sk,c) \rightarrow m'$

 $Evaluate_{\epsilon}(f,c_{1},...,c_{i}) \rightarrow c$

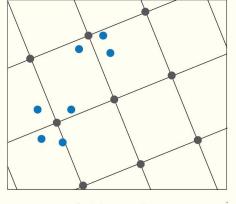
The secret key, sk, is a random P-bit odd integer p.

A somewhat homomorphic encryption scheme

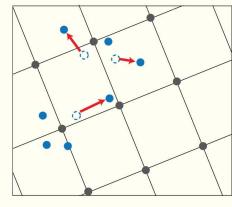
- A limited number of functions supported.

Lattice-based cryptography

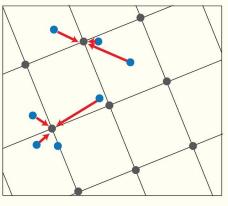
Lattice-based cryptography (for fully homomorphic encryption)



1. Encrypt



2. Add noise





Bootstrapping

noise =
$$p/2$$

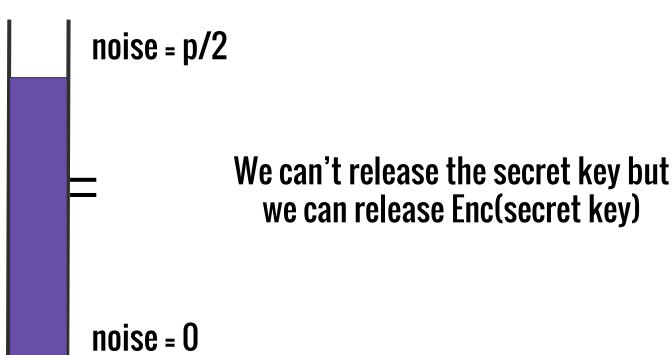
Addition doubles noise Multiplication squares

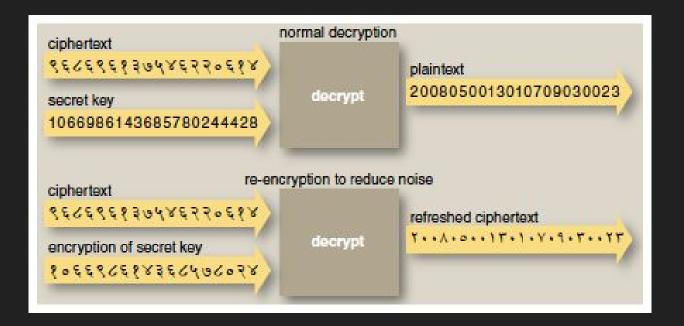
noise = O

Bootstrapping



Bootstrapping

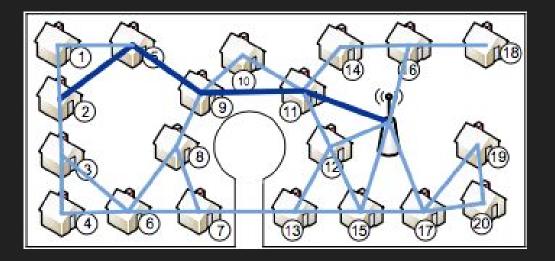




Is Gentry's scheme practical?

Implementing homomorphic encryption in the smart grid

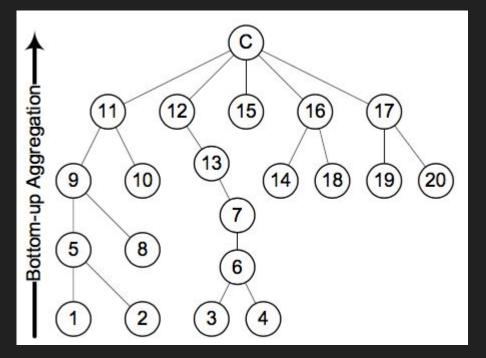
An example of smart grid communication in a neighborhood.



[-] Excessive network traffic

[-] High overhead for collector

Improved communication graph

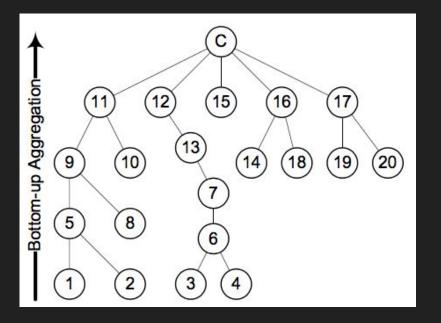


Collector calculates a BFS Tree [+] Collector offloads computation [+] Computation is done in a distributed manner

+] Tree can easily be rebalanced

How is homomorphic encryption used in this scheme?

System design



7-tuple message

{TID, Trigger, Data, Collect, Operation, Destination, Key}

Child node

Generate its readings, encrypts data and sends to parent

Parent node

Collects data from all children, appends its own reading Sends data to its parent or collector

Collector node

- Collects data from children
- Performs some operation
- Decrypts data

Paillier Cryptosystem - Key Generation

Key Generation

- 1. Pick two large prime numbers p and q;
- 2. N = p \cdot q and λ = lcm(p-1, q-1), where lcm represents least common multiple.
- 3. Select a random number g where g \in $Z^*_{n^2}$
- 4. Set function L(u) as: L(u) = (u 1)/N

- 5. Ensure that N divides the order of g: check if $L(g^{\lambda} \mod N^2)$ and n are co-prime, i.e. $gcd(L(g^{\lambda} \mod N^2), N)=1$
- 6. (N, g) is the public key pair.
- 7. (p, q) is the private key pair.

Paillier Cryptosystem - Encryption

- 1. We want to encrypt the message: $m \in Z_{N}^{*}$
- 2. Select a random number: $r \in Z_{N}^{*}$
- 3. Encrypt m using: $c = E(m) = g^{m} \cdot r^{N} \mod N^{2}$

Paillier Cryptosystem - Decryption

1. We want to decrypt ciphertext: $c \in Z^*_{N^2}$

2) Decrypt with: $m = \mathsf{D}(c) = \left(\frac{L(c^{\lambda} \mod N^2)}{L(g^{\lambda} \mod N^2)}\right) \mod N$

Given
$$c_1 = E(m_1)$$
 and $c_2 = E(m_2)$, $\forall m_1, m_2 \in Z_N$, we have $D(c_1 \cdot c_2 \mod N^2) = m_1 + m_2 \mod N$

Medical Applications

Financial Applications

Advertising and Pricing

Voting

Summary

- Properties of homomorphic encryption
- Partial homomorphism
- Fully homomorphism
- An example how it can be integrated into the smart grid